

The large deviation principle for the Erdős-Rényi random graph

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Main objective: how to count graphs with a given property

- ▶ Only consider finite undirected graphs without self-loops in this talk.
- ▶ $2^{\binom{n-1}{2}}$ such graphs on n vertices.
- ▶ Question: Given a property P and an integer n , roughly **how many of these graphs have property P ?**
- ▶ For example, P may be: $\#\text{triangles} \geq tn^3$, where t is a given constant.
- ▶ To make any progress, need to assume some regularity on P . For example, we may demand that P be **continuous with respect to some metric**.
- ▶ **What metric? What space?**

Another motivation

- ▶ Let $G(n, p)$ be the Erdős-Rényi random graph on n vertices where each edge is added independently with probability p .
- ▶ Number of triangles in $G(n, p)$ roughly $\binom{n}{3}p^3 \sim n^3p^3/6$.
- ▶ What if, just by chance, #triangles turns out to be $\approx tn^3$ where $t > p^3/6$? What would the graph look like, conditional on this rare event?

An abstract topological space of graphs

- ▶ Beautiful unifying theory developed by Lovász and coauthors (listed in order of frequency: V. T. Sós, B. Szegedy, C. Borgs, J. Chayes, K. Vesztegombi, A. Schrijver and M. Freedman). Related to earlier works of Aldous, Hoover, Kallenberg.
- ▶ Let G_n be a sequence of simple graphs whose number of nodes tends to infinity.
- ▶ For every fixed simple graph H , let $|\text{hom}(H, G)|$ denote the number of homomorphisms of H into G (i.e. edge-preserving maps $V(H) \rightarrow V(G)$, where $V(H)$ and $V(G)$ are the vertex sets).
- ▶ This number is normalized to get the **homomorphism density**

$$t(H, G) := \frac{|\text{hom}(H, G)|}{|V(G)|^{|V(H)|}.$$

This gives the probability that a random mapping $V(H) \rightarrow V(G)$ is a homomorphism.

Abstract space of graphs contd.

- ▶ Suppose that $t(H, G_n)$ tends to a limit $t(H)$ for every H .
- ▶ Then Lovász & Szegedy proved that there is a natural “limit object” in the form of a function $f \in \mathcal{W}$, where \mathcal{W} is the space of all measurable functions from $[0, 1]^2$ into $[0, 1]$ that satisfy $f(x, y) = f(y, x)$ for all x, y .
- ▶ Conversely, every such function arises as the limit of an appropriate graph sequence.
- ▶ This limit object determines all the limits of subgraph densities: if H is a simple graph with k vertices, then

$$t(H, f) = \int_{[0,1]^k} \prod_{(i,j) \in E(H)} f(x_i, x_j) dx_1 \cdots dx_k.$$

- ▶ A sequence of graphs $\{G_n\}_{n \geq 1}$ is said to converge to f if for every finite simple graph H ,

$$\lim_{n \rightarrow \infty} t(H, G_n) = t(H, f).$$

Example

- ▶ For any fixed graph H ,

$$t(H, G(n, p)) \rightarrow p^{|E(H)|} \text{ almost surely as } n \rightarrow \infty.$$

- ▶ On the other hand, if f is the function that is identically equal to p , then $t(H, f) = p^{|E(H)|}$.
- ▶ Thus, the sequence of random graphs $G(n, p)$ converges almost surely to the non-random limit function $f(x, y) \equiv p$ as $n \rightarrow \infty$.

Abstract space of graphs contd.

- ▶ The elements of \mathcal{W} are sometimes called 'graphons'.
- ▶ A finite simple graph G on n vertices can also be represented as a graphon f^G in a natural way:

$$f^G(x, y) = \begin{cases} 1 & \text{if } (\lceil nx \rceil, \lceil ny \rceil) \text{ is an edge in } G, \\ 0 & \text{otherwise.} \end{cases}$$

- ▶ Note that this allows *all* simple graphs, irrespective of the number of vertices, to be represented as elements of the single abstract space \mathcal{W} .
- ▶ So, what is the topology on this space?

The cut metric

- ▶ For any $f, g \in \mathcal{W}$, Frieze and Kannan defined the cut distance:

$$d_{\square}(f, g) := \sup_{S, T \subseteq [0,1]} \left| \int_{S \times T} [f(x, y) - g(x, y)] dx dy \right|.$$

- ▶ Introduce an equivalence relation on \mathcal{W} : say that $f \sim g$ if $f(x, y) = g_{\sigma}(x, y) := g(\sigma x, \sigma y)$ for some measure preserving bijection σ of $[0, 1]$.
- ▶ Denote by \tilde{g} the closure in $(\mathcal{W}, d_{\square})$ of the orbit $\{g_{\sigma}\}$.
- ▶ The quotient space is denoted by $\widetilde{\mathcal{W}}$ and τ denotes the natural map $g \rightarrow \tilde{g}$.
- ▶ Since d_{\square} is invariant under σ one can define on $\widetilde{\mathcal{W}}$ the natural distance δ_{\square} by

$$\delta_{\square}(\tilde{f}, \tilde{g}) := \inf_{\sigma} d_{\square}(f, g_{\sigma}) = \inf_{\sigma} d_{\square}(f_{\sigma}, g) = \inf_{\sigma_1, \sigma_2} d_{\square}(f_{\sigma_1}, g_{\sigma_2})$$

making $(\widetilde{\mathcal{W}}, \delta_{\square})$ into a metric space.

Cut metric and graph limits

To any finite graph G , we associate the natural graphon f^G and its orbit $\tilde{G} = \tau f^G = \tilde{f}^G \in \tilde{\mathcal{W}}$. One of the key results of the is the following:

Theorem (Borgs, Chayes, Lovász, Sós & Vesztergombi)

A sequence of graphs $\{G_n\}_{n \geq 1}$ converges to a limit $f \in \mathcal{W}$ if and only if $\delta_{\square}(\tilde{G}_n, \tilde{f}) \rightarrow 0$ as $n \rightarrow \infty$.

Our result

- ▶ For any Borel set $\tilde{A} \subseteq \tilde{\mathcal{W}}$, let

$$\tilde{A}_n := \{\tilde{h} \in \tilde{A} : \tilde{h} = \tilde{G} \text{ for some } G \text{ on } n \text{ vertices}\}.$$

- ▶ Let $I(u) := \frac{1}{2}u \log u + \frac{1}{2}(1-u) \log(1-u)$.
- ▶ For any $\tilde{h} \in \tilde{\mathcal{W}}$, let $I(\tilde{h}) := \iint I(h(x,y)) dx dy$, where h is any element of \tilde{h} .

Theorem (Chatterjee & Varadhan, 2010)

The function I is well-defined and lower-semicontinuous on $\tilde{\mathcal{W}}$. If \tilde{F} is a closed subset of $\tilde{\mathcal{W}}$ then

$$\limsup_{n \rightarrow \infty} n^{-2} \log |\tilde{F}_n| \leq - \inf_{\tilde{h} \in \tilde{F}} I(\tilde{h})$$

and if \tilde{U} is an open subset of $\tilde{\mathcal{W}}$, then

$$\liminf_{n \rightarrow \infty} n^{-2} \log |\tilde{U}_n| \geq - \inf_{\tilde{h} \in \tilde{U}} I(\tilde{h}).$$

- ▶ Counting graphs can be related to finding large deviation probabilities for Erdős-Rényi random graphs.
- ▶ For example,

$$\begin{aligned} & \# \text{graphs on } n \text{ vertices satisfying } P \\ &= 2^{n(n-1)/2} \mathbb{P}(G(n, 1/2) \text{ satisfies } P). \end{aligned}$$

- ▶ Indeed, the main result in our paper is stated as a large deviation principle for the Erdős-Rényi graph, which can be easily proved to be equivalent to the graph counting principle stated before.

Large deviation principle for ER graphs

- ▶ The random graph $G(n, p)$ induces probability distribution $\tilde{\mathbb{P}}_{n,p}$ on the space $\tilde{\mathcal{W}}$ through the map $G \rightarrow \tilde{G}$.
- ▶ Let $I_p(u) := \frac{1}{2}u \log \frac{u}{p} + \frac{1}{2}(1-u) \log \frac{1-u}{1-p}$.
- ▶ For $\tilde{h} \in \tilde{\mathcal{W}}$, let $I_p(\tilde{h}) := \iint I_p(h(x, y)) dx dy$, where h is any element of \tilde{h} .

Theorem (Chatterjee & Varadhan, 2010)

For any closed set $\tilde{F} \subseteq \tilde{\mathcal{W}}$,

$$\limsup_{n \rightarrow \infty} \frac{1}{n^2} \log \tilde{\mathbb{P}}_{n,p}(\tilde{F}) \leq - \inf_{\tilde{h} \in \tilde{F}} I_p(\tilde{h}).$$

and for any open set $\tilde{U} \subseteq \tilde{\mathcal{W}}$,

$$\liminf_{n \rightarrow \infty} \frac{1}{n^2} \log \tilde{\mathbb{P}}_{n,p}(\tilde{U}) \geq - \inf_{\tilde{h} \in \tilde{U}} I_p(\tilde{h}).$$

- ▶ The LDP can be proved by standard techniques for the weak topology on $\widetilde{\mathcal{W}}$. (Fenchel-Legendre transforms, Gärtner-Ellis theorem, etc.)
- ▶ However, the weak topology is not very interesting. For example, subgraph counts are not continuous with respect to the weak topology.
- ▶ The LDP for the topology of the cut metric does not follow via standard methods.

Szemerédi's lemma

- ▶ Let $G = (V, E)$ be a simple graph of order n .
- ▶ For any $X, Y \subseteq V$, let $e_G(X, Y)$ be the number of X - Y edges of G and let

$$\rho_G(X, Y) := \frac{e_G(X, Y)}{|X||Y|}$$

- ▶ Call a pair (A, B) of disjoint sets $A, B \subseteq V$ **ϵ -regular** if all $X \subseteq A$ and $Y \subseteq B$ with $|X| \geq \epsilon|A|$ and $|Y| \geq \epsilon|B|$ satisfy $|\rho_G(X, Y) - \rho_G(A, B)| \leq \epsilon$.
- ▶ A partition $\{V_0, \dots, V_K\}$ of V is called an **ϵ -regular partition of G** if it satisfies the following conditions: (i) $|V_0| \leq \epsilon n$; (ii) $|V_1| = |V_2| = \dots = |V_K|$; (iii) all but at most ϵK^2 of the pairs (V_i, V_j) with $1 \leq i < j \leq K$ are ϵ -regular.

Theorem (Szemerédi's lemma)

Given $\epsilon > 0$, $m \geq 1$ there exists $M = M(\epsilon, m)$ such that every graph of order $\geq M$ admits an ϵ -regular partition $\{V_0, \dots, V_K\}$ for some $K \in [m, M]$.

Finishing the proof using Szemerédi's lemma

- ▶ Suppose G is a graph of order n with ϵ -regular partition $\{V_0, \dots, V_K\}$.
- ▶ Let G' be the random graph with independent edges where a vertex $u \in V_i$ is connected to a vertex $v \in V_j$ with probability $\rho_G(V_i, V_j)$.
- ▶ Using Szemerédi's regularity lemma, one can prove that $\delta_{\square}(G, G') \simeq 0$ with high probability if K and n are appropriately large and ϵ is small.
- ▶ The chance that $G(n, p) \approx G'$ is computed by approximating the probability density of G' with respect to $G(n, p)$ as in the proof of Cramér's theorem.
- ▶ Combining the two steps gives the probability that $G(n, p) \approx G$, leading to the LDP for $G(n, p)$, and finally leading to the graph counting theorem.

Conditional distributions

Theorem

Take any $p \in (0, 1)$. Let \tilde{F} be a closed subset of $\tilde{\mathcal{W}}$ satisfying

$$\inf_{\tilde{h} \in \tilde{F}^o} I_p(\tilde{h}) = \inf_{\tilde{h} \in \tilde{F}} I_p(\tilde{h}) > 0.$$

Let \tilde{F}^* be the subset of \tilde{F} where I_p is minimized. Then \tilde{F}^* is *non-empty and compact*, and for each n , and each $\epsilon > 0$,

$$\mathbb{P}(\delta_{\square}(G(n, p), \tilde{F}^*) \geq \epsilon \mid G(n, p) \in \tilde{F}) \leq e^{-C(\epsilon, \tilde{F})n^2}$$

where $C(\epsilon, \tilde{F})$ is a positive constant depending only on ϵ and \tilde{F} .

Proof: Follows from the compactness of $\tilde{\mathcal{W}}$ (a deep result of Lovász and Szegedy, involving recursive applications of Szemerédi's lemma and martingales).

Large deviations for triangle counts

- ▶ Let $T_{n,p}$ be the number of triangles in $G(n, p)$.
- ▶ Objective: to evaluate the limit

$$\lim_{n \rightarrow \infty} \frac{1}{n^2} \log \mathbb{P}(T_{n,p} \geq (1 + \epsilon) \mathbb{E}(T_{n,p}))$$

as a function of p and ϵ .

- ▶ In the sparse case ($p \rightarrow 0$), a long-standing conjecture about matching the upper and lower bounds was recently resolved by [Chatterjee](#) and also shortly afterward by [DeMarco & Kahn](#).
- ▶ For fixed p , exact evaluation of limit due to [Chatterjee & Dey](#): for a certain explicit set of (p, t) ,

$$\lim_{n \rightarrow \infty} \frac{1}{n^2} \log \mathbb{P}(T_{n,p} \geq tn^3) = -I_p((6t)^{1/3}),$$

when $I_p(u) := \frac{1}{2}u \log \frac{u}{p} + \frac{1}{2}(1-u) \log \frac{1-u}{1-p}$.

- ▶ Unfortunately, the result does not cover all values of (p, t) .

Large deviations for triangle counts contd.

- ▶ Recall: \mathcal{W} is the space of symmetric measurable functions from $[0, 1]^2$ into $[0, 1]$.
- ▶ For each $f \in \mathcal{W}$, let

$$T(f) := \frac{1}{6} \int_0^1 \int_0^1 \int_0^1 f(x, y) f(y, z) f(z, x) dx dy dz$$

and let $I_p(f) = \iint I_p(f(x, y)) dx dy$.

- ▶ For each $p \in (0, 1)$ and $t \geq 0$, let

$$\phi(p, t) := \inf \{ I_p(f) : f \in \mathcal{W}, T(f) \geq t \}. \quad (1)$$

Theorem (Chatterjee & Varadhan, 2010)

For each $p \in (0, 1)$ and each $t \geq 0$,

$$\lim_{n \rightarrow \infty} \frac{1}{n^2} \log \mathbb{P}(T_{n,p} \geq tn^3) = -\phi(p, t).$$

Moreover, the infimum is attained in the variational problem (1).

The 'replica symmetric' phase

Theorem (Chatterjee & Varadhan, 2010)

Let $h_p(t) := I_p((6t)^{1/3})$. Let \hat{h}_p be the convex minorant of h_p . If t is a point where $h_p(t) = \hat{h}_p(t)$, then $\phi(p, t) = h_p(t)$. Moreover, for such (p, t) , the conditional distribution of $G(n, p)$ given $T_{n,p} \geq tn^3$ is indistinguishable from the law of $G(n, (6t)^{1/3})$ in the large n limit.

Remarks: This result recovers the result of Chatterjee & Dey and gives more. However, the theorem of Chatterjee & Dey gives an error bound of order $n^{-1/2}$, which is impossible to obtain via Szemerédi's lemma.

'Replica symmetry breaking'

The following theorem shows that given any t , for all p small enough, the conditional distribution of $G(n, p)$ given $T_{n,p} \geq tn^3$ **does not** resemble that of an Erdős-Rényi graph.

Theorem (Chatterjee & Varadhan, 2010)

Let $\tilde{\mathcal{C}}$ denote the set of constant functions in $\tilde{\mathcal{W}}$ (representing all Erdős-Rényi graphs). For each t , there exists $p' > 0$ and $\epsilon > 0$ such that for all $p < p'$,

$$\lim_{n \rightarrow \infty} \mathbb{P}(\delta_{\square}(G(n, p), \tilde{\mathcal{C}}) > \epsilon \mid T_{n,p} \geq tn^3) = 1.$$

The double phase transition

Theorem (Chatterjee & Varadhan, 2010)

There exists $p_0 > 0$ such that if $p \leq p_0$, then there exists $p^3/6 < t' < t'' < 1/6$ such that the replica symmetric picture holds when $t \in (p^3/6, t') \cup (t'', 1/6)$, but there is a non-empty subset of (t', t'') where replica symmetry breaks down.

The small p limit

The following theorem says that when t is fixed and p is very small, then conditionally on the event $\{T_{n,p} \geq tn^3\}$ the graph $G(n, p)$ must look like a clique.

Theorem (Chatterjee & Varadhan, 2010)

For each t ,

$$\lim_{p \rightarrow 0} \frac{\phi(p, t)}{\log(1/p)} = \frac{(6t)^{2/3}}{2}.$$

Moreover, if

$$\chi_t(x, y) := \mathbf{1}_{\{\max\{x, y\} \leq (6t)^{1/3}\}}$$

is the graphon representing a clique with triangle density t , then for each $\epsilon > 0$,

$$\lim_{p \rightarrow 0} \lim_{n \rightarrow \infty} \mathbb{P}(\delta_{\square}(\widetilde{G}(n, p), \widetilde{\chi}_t) \geq \epsilon \mid T_{n,p} \geq tn^3) = 0.$$

Lower tails

- ▶ Given a fixed simple graph H ,

$$\lim_{u \rightarrow 0} \lim_{n \rightarrow \infty} \frac{\log \mathbb{P}(t(H, G(n, p)) \leq u)}{n^2} = -\frac{1}{2(\chi(H) - 1)} \log \frac{1}{1 - p},$$

where $\chi(H)$ is the chromatic number of H .

- ▶ Closely related to the Erdős-Stone theorem from extremal graph theory.
- ▶ In fact, the precise result implies the following: given that $t(H, G(n, p))$ is very small (or zero), the graph $G(n, p)$ looks like a complete $(\chi(H) - 1)$ -partite graph with $(1 - p)$ -fraction of edges randomly deleted.
- ▶ However, if $t(H, G(n, p))$ is just a little bit below its expected value, the graph continues to look like an Erdős-Rényi graph as in the upper tail case.

Open questions

- ▶ There are many questions that remain unresolved, even in the simple example of upper tails for triangle counts. For example:
- ▶ What is the set of optimal solutions of the variational problem defining the rate function in the broken replica symmetry phase (i.e. where the optimizer is not a constant)?
- ▶ Is the solution unique in the quotient space $\widetilde{\mathcal{W}}$, or can there exist multiple solutions?
- ▶ Is it possible to explicitly compute a nontrivial solution for at least some values of (p, t) in the broken replica symmetry region?
- ▶ Is it possible to even numerically evaluate or approximate a solution using a computer?
- ▶ What is the full characterization of the replica symmetric phase? What is the phase boundary?
- ▶ What happens in the sparse case where p and t are both allowed to tend to zero?

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