

Preliminary Exam

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March 17th, 2020

- Study self-duality of complete intersections
- Scheja-Storch: construct explicit isomorphism Θ for self-duality
- **Goal:** describe Scheja-Storch's isomorphism
- **Context:** A a commutative ring, B a finitely generated projective A -algebra

- A a commutative ring, B a finitely generated projective A -algebra
- $\chi : B \otimes_A B \rightarrow \text{Hom}_A(\text{Hom}_A(B, A), B)$ defined by
$$\chi(b \otimes c) = (\varphi \mapsto \varphi(b)c)$$
- Presentation of B as a complete intersection over A :
 - $A[x]$ denotes either $A[x_1, \dots, x_n]$ or $A[[x_1, \dots, x_n]]$
 - $\rho : A[x] \rightarrow B$ with $\ker \rho = (t_1, \dots, t_n)$
- Kähler differentials $dt_j = t_j \otimes 1 - 1 \otimes t_j = \sum_{i=1}^n a_{ij}(x_i \otimes 1 - 1 \otimes x_i)$
 - $\mu_{A[x]} : A[x] \otimes_A A[x] \rightarrow A[x]$ given by $b \otimes c \mapsto bc$
 - $\ker \mu_{A[x]} = (x_i \otimes 1 - 1 \otimes x_i : 1 \leq i \leq n)$
 - $t_j \otimes 1 - 1 \otimes t_j \in \ker \mu_{A[x]}$
- $\Delta = \rho \otimes \rho(\det(a_{ij})) \in B \otimes_A B$
 - **Remark:** $\mu_B(\Delta) = \rho(\det(\partial t_j / \partial x_i)) := J$

- Scheja-Storch: construct explicit isomorphism between B and B^*

Theorem

The homomorphism $\Theta := \chi(\Delta) : \text{Hom}_A(B, A) \rightarrow B$ is a B -linear isomorphism.

- Proof uses the following lemmas:
 - **Lemma 1:** χ induces a B -linear isomorphism $\text{Ann}_{B \otimes_A B}(\ker \mu_B) \rightarrow \text{Hom}_B(\text{Hom}_A(B, A), B)$.
 - **Lemma 2:** $(B \otimes_A B) \cdot \Delta = \text{Ann}_{B \otimes_A B}(\ker \mu_B)$ and $\ker \mu_B = \text{Ann}_{B \otimes_A B}((B \otimes_A B) \cdot \Delta)$.
 - **Lemma 3:** $\text{Hom}_A(B, A)$ is a projective A -module.
 - **Lemma 4:** $\text{Hom}_A(B, A)$ is a free B -module.

Proof of Main Theorem

Theorem: $\Theta := \chi(\Delta) : \text{Hom}_A(B, A) \rightarrow B$ is a B -linear isomorphism.

- $(B \otimes_A B) \cdot \Delta = ((B \otimes_A B) / \text{Ann}_{B \otimes_A B} \Delta) \cdot \Delta$
 - **Lemma 2** $\implies \text{Ann}_{B \otimes_A B} \Delta = \ker \mu_B$
 - $(B \otimes_A B) / \text{Ann}_{B \otimes_A B} \Delta = (B \otimes_A B) / \ker \mu_B$
 - **First iso. theorem for modules** $\implies (B \otimes_A B) / \ker \mu_B \cong B$
 - Thus $(B \otimes_A B) \cdot \Delta \cong B \cdot \Delta$
- **Lemma 2** $\implies \text{Ann}_{B \otimes_A B} \ker \mu_B = (B \otimes_A B) \cdot \Delta = B \cdot \Delta$
 - $\text{Ann}_{B \otimes_A B} \ker \mu_B$ is a free B -module with basis Δ
 - **Lemma 1** $\implies \chi(\text{Ann}_{B \otimes_A B} \ker \mu_B)$ is a free B -module with basis $\chi(\Delta)$
 - **Lemma 1** $\implies \text{Hom}_B(\text{Hom}_A(B, A), B)$ is a free B -module with basis $\chi(\Delta)$
- **Lemmas 3,4** $\implies \text{Hom}_A(B, A)$ is a free B -module of rank 1
 - $\chi(\Delta) : \text{Hom}_A(B, A) \rightarrow B$ is B -linear homomorphism of rank 1 free B -modules
 - Any $\text{Hom}_A(B, A) \rightarrow B$ is a B -multiple of $\chi(\Delta)$, so $\chi(\Delta)$ is an isomorphism

Proof of Lemma 1

Lemma 1: χ induces $\text{Ann}_{B \otimes_A B} \ker \mu_B \cong \text{Hom}_B(\text{Hom}_A(B, A), B)$.

- $\chi : B \otimes_A B \rightarrow \text{Hom}_A(\text{Hom}_A(B, A), B)$ is a bijection
 - Look at two B -module structures compatible with χ
- On $B \otimes_A B$, let $a \cdot (b \otimes c) = ab \otimes c$ and $a * (b \otimes c) = b \otimes ac$
 - $\ker \mu_B = (a \otimes 1 - 1 \otimes a : a \in B)$
 - $\text{Ann}_{B \otimes_A B} \ker \mu_B = \{x \in B \otimes_A B : (a \otimes 1 - 1 \otimes a)x = 0 \forall a \in B\}$
 - $\text{Ann}_{B \otimes_A B} \ker \mu_B = \{x \in B \otimes_A B : a \cdot x = a * x \forall a \in B\}$
 - $\implies \text{Ann}_{B \otimes_A B} \ker \mu_B$ is the largest submodule where \cdot and $*$ agree
- On $\text{Hom}_A(\text{Hom}_A(B, A), B)$, let $a \cdot \varphi = (\psi \mapsto \varphi(a\psi))$ and $a * \varphi = (\psi \mapsto a\varphi(\psi))$
 - $a \cdot \varphi = a * \varphi$ for all $a \in B$ iff φ is B -linear
 - $\implies \text{Hom}_B(\text{Hom}_A(B, A), B)$ is the largest submodule where \cdot and $*$ agree
- $\chi(a \cdot (b \otimes c)) = a \cdot \chi(b \otimes c)$ and $\chi(a * (b \otimes c)) = a * \chi(b \otimes c)$

Proof of Lemma 2

Lemma 2: $(B \otimes_A B) \cdot \Delta = \text{Ann}_{B \otimes_A B} \ker \mu_B$ and
 $\ker \mu_B = \text{Ann}_{B \otimes_A B}((B \otimes_A B) \cdot \Delta)$

- Reduce to the following lemma:

Lemma 5: Let A be a commutative ring. Assume:

- Ideals $(g_1, \dots, g_n) = \mathfrak{b} \subseteq \mathfrak{a} = (f_1, \dots, f_n)$
- $g_j = \sum_{i=1}^n a_{ij} f_i$ and $\Delta = \det(a_{ij})$
- $\mathfrak{a}' = \mathfrak{a}/\mathfrak{b}$ in $A' = A/\mathfrak{b}$ and Δ' image of Δ in A'
- f_1, \dots, f_n and g_1, \dots, g_n are prime sequences in $A_{\mathfrak{p}}$ for all $\mathfrak{p} \supseteq \mathfrak{a}$

Then:

- Δ' is independent of choice of a_{ij}
- $\Delta' A' = \text{Fitt}_{A'}(\mathfrak{a}')$
- $\Delta' A' = \text{Ann}_{A'}(\mathfrak{a}')$ and $\mathfrak{a}' = \text{Ann}_{A'}(\Delta' A')$

Proof of Lemma 2

Lemma 2: $(B \otimes_A B) \cdot \Delta = \text{Ann}_{B \otimes_A B} \ker \mu_B$ and
 $\ker \mu_B = \text{Ann}_{B \otimes_A B}((B \otimes_A B) \cdot \Delta)$

$$\begin{array}{ccc}
 A[x] \otimes_A A[x] & \xrightarrow{\mu_{A[x]}} & A[x] \\
 \downarrow \rho \otimes \text{id} & & \downarrow \rho \\
 B \otimes_A A[x] & & B \\
 \downarrow \text{id} \otimes \rho & & \\
 B \otimes_A B & \xrightarrow{\mu_B} & B
 \end{array}$$

- For $A[x] = A[[x_1, \dots, x_n]]$, consider $\text{id} \otimes \rho : B \otimes_A A[x] \rightarrow B \otimes_A B$
 - $\ker(\text{id} \otimes \rho) = (-1 \otimes t_1, \dots, -1 \otimes t_n)$ plays the role of \mathfrak{b}
 - $(\rho(x_1) \otimes 1 - 1 \otimes x_1, \dots, \rho(x_n) \otimes 1 - 1 \otimes x_n)$ plays the role of \mathfrak{a}
- For $A[x] = A[x_1, \dots, x_n]$, reduce to the local case in two steps:
 - Locally in B , i.e. over $B_{\mathfrak{n}} \otimes_A B_{\mathfrak{n}}$
 - Locally in $B \otimes_A B$, i.e. over $(B \otimes_A B)_{\mathfrak{m}}$
 - Clear denominators to go from locally in B to locally in $B \otimes_A B$

Proof of Lemmas 3 and 4

Lemma 3: $\text{Hom}_A(B, A)$ is a projective A -module.

- B finitely generated $\implies \pi : A^m \twoheadrightarrow B$
 - $0 \rightarrow \ker \pi \rightarrow A^m \rightarrow B \rightarrow 0$ and B projective $\implies B \oplus \ker \pi \cong A^m$
 - $A^m \cong \text{Hom}_A(A^m, A) \cong \text{Hom}_A(B \oplus \ker \pi, A) \cong \text{Hom}_A(B, A) \oplus \text{Hom}_A(\ker \pi, A)$

Lemma 4: $\text{Hom}_A(B, A)$ is a free B -module.

- If M finitely generated B -module, M projective A -module, and $\text{Hom}_B(M, B)$ projective (free) B -module, then M is a projective (free) B -module.

Proof of Lemma 5

Lemma 5: Let A be a commutative ring. Assume:

- Ideals $(g_1, \dots, g_n) = \mathfrak{b} \subseteq \mathfrak{a} = (f_1, \dots, f_n)$
- $g_j = \sum_{i=1}^n a_{ij} f_i$ and $\Delta = \det(a_{ij})$
- $\mathfrak{a}' = \mathfrak{a}/\mathfrak{b}$ in A/\mathfrak{b} and Δ' image of Δ in A'
- f_1, \dots, f_n and g_1, \dots, g_n are prime sequences in $A_{\mathfrak{p}}$ for all $\mathfrak{p} \supseteq \mathfrak{a}$

Then:

- (a) Δ' is independent of choice of a_{ij}
- (b) $\Delta' A' = \text{Fitt}_{A'}(\mathfrak{a}')$
- (c) $\Delta' A' = \text{Ann}_{A'}(\mathfrak{a}')$ and $\mathfrak{a}' = \text{Ann}_{A'}(\Delta' A')$
 - A *prime sequence* is a regular sequence f_1, \dots, f_n such that (f_1, \dots, f_i) is a prime ideal for each $1 \leq i \leq n$
 - The *Fitting ideal* $\text{Fitt}_{A'}(\mathfrak{a}')$ is generated by the determinants of all $n \times n$ minors $\det(b_{ij})$ of the collection of syzygies $\sum_{i=1}^n b_{ij} f_i = 0$, where f_1, \dots, f_n generate \mathfrak{a}' over A'

Proof of Lemma 5

Lemma 5: Let A be a commutative ring. Assume:

- Ideals $(g_1, \dots, g_n) = \mathfrak{b} \subseteq \mathfrak{a} = (f_1, \dots, f_n)$
- $g_j = \sum_{i=1}^n a_{ij} f_i$ and $\Delta = \det(a_{ij})$
- $\mathfrak{a}' = \mathfrak{a}/\mathfrak{b}$ in A/\mathfrak{b} and Δ' image of Δ in A'
- f_1, \dots, f_n and g_1, \dots, g_n are prime sequences in $A_{\mathfrak{p}}$ for all $\mathfrak{p} \supseteq \mathfrak{a}$

Then:

- (a) Δ' is independent of choice of a_{ij}
- (b) $\Delta' A' = \text{Fitt}_{A'}(\mathfrak{a}')$
- (c) $\Delta' A' = \text{Ann}_{A'}(\mathfrak{a}')$ and $\mathfrak{a}' = \text{Ann}_{A'}(\Delta' A')$
 - (a), (b), (c) are local in A'
 - If $\mathfrak{a}' = A'$, then \mathfrak{a}' is one-dimensional over A'
 - Any $n \times n$ minor contains linearly dependent rows, so $\Delta' = 0$
 - May assume A' is local, $\{f_i\}$ and $\{g_i\}$ regular sequences contained in unique max ideal

Proof of Lemma 5

(a) Δ' is independent of choice of a_{ij} .

- Suppose $g_j = \sum_{i=1}^n a_{ij} f_i = \sum_{i=1}^n b_{ij} f_i$
 - Need to show $\det(a_{ij}) - \det(b_{ij}) \in \mathfrak{b}$
 - WLOG, assume first $n - 1$ rows of (a_{ij}) and (b_{ij}) agree
 - Set $c_{ij} = a_{ij} = b_{ij}$ for $1 \leq j \leq n - 1$ and $c_{in} = a_{in} - b_{in}$
 - $\det(c_{ij}) = \det(a_{ij}) - \det(b_{ij})$
- Cramer's rule \implies
 $\det(c_{ij}) \cdot f_i = \det(\text{replace } i^{\text{th}} \text{ column of } (c_{ij}) \text{ with } (g_1, \dots, g_{n-1}, 0)^T)$
 - $\implies \det(c_{ij}) \cdot \mathfrak{a} \subseteq (g_1, \dots, g_{n-1})$
 - g_1, \dots, g_n prime sequence $\implies (g_1, \dots, g_{n-1})$ a prime ideal
 - g_1, \dots, g_n a regular sequence in $\mathfrak{a} \implies \mathfrak{a} \not\subseteq (g_1, \dots, g_{n-1})$
 - $\implies \det(c_{ij}) \in (g_1, \dots, g_{n-1}) \subseteq \mathfrak{b}$

(b) $\Delta' A' = \text{Fitt}_{A'}(\mathfrak{a}')$.

- $\text{Fitt}_{A'}(\mathfrak{a}')$ is the image of $\text{Fitt}_A(\mathfrak{a})$ in A'
- $\mathfrak{a}' \cong A^n/U$
 - U is generated by (a_{1j}, \dots, a_{nj}) and $(0, \dots, f_q, \dots, -f_p, \dots, 0)$
 - One of these minors is Δ
 - Using Cramer's rule as in (a), all other minors are in \mathfrak{b}
- Fitting ideal does not depend on choice of generators or relations

Proof of Lemma 5

(c) $\Delta' A' = \text{Ann}_{A'}(\mathfrak{a}')$ and $\mathfrak{a}' = \text{Ann}_{A'}(\Delta' A')$.

- Induct on n
- $n = 0$:
 - $\Delta' = \det(\emptyset) = \prod_{\emptyset} \lambda = 1$ and $\mathfrak{a}' = 0$
 - Trivially, $\text{Ann}_{A'}(0) = A'$ and $0 = \text{Ann}_{A'}(A')$
- $n \geq 1$:
 - (b) $\implies \Delta' A'$ is a Fitting ideal
 - $\implies \Delta' A'$ does not depend on choice of f_1, \dots, f_n
 - Pick f_1, \dots, f_n so that f_1 is prime to (g_2, \dots, g_n)
 - Set $B = A/(g_2, \dots, g_n)$, $b_{ij} = a_{ij}$ for $j \geq 2$, $b_{11} = 1$ and $b_{i1} = 0$ for $i \geq 2$, $\Delta_b = \det(b_{ij})$
 - Inductive hypothesis on f_2, \dots, f_n and g_2, \dots, g_n in $A/f_1 A \implies$:
 - (i) $(B/f_1 B) \cdot \Delta_b = \text{Ann}_{B/f_1 B}(\mathfrak{a}B/f_1 B) = (f_1 B : \mathfrak{a}B)/f_1 B$
 - (ii) $\mathfrak{a}B/f_1 B = \text{Ann}_{B/f_1 B}((B/f_1 B) \cdot \Delta_b) = (f_1 B : \Delta_b B)/f_1 B$
 - (ii) $\implies \mathfrak{a}B = (f_1 B : \Delta_b B)$
 - Cramer's rule $\implies \Delta f_1, \Delta_b g_1 \in \mathfrak{b} \implies \Delta_b g_1 \equiv \Delta f_1 \pmod{\mathfrak{b}}$

Proof of Lemma 5

(c) $\Delta' A' = \text{Ann}_{A'}(\mathfrak{a}')$ and $\mathfrak{a}' = \text{Ann}_{A'}(\Delta' A')$.

- Inductive hypothesis on f_2, \dots, f_n and g_2, \dots, g_n in $A/f_1 A \implies$:
 - (i) $(B/f_1 B) \cdot \Delta_b = \text{Ann}_{B/f_1 B}(\mathfrak{a}B/f_1 B) = (f_1 B : \mathfrak{a}B)/f_1 B$
 - (ii) $\mathfrak{a}B/f_1 B = \text{Ann}_{B/f_1 B}((B/f_1 B) \cdot \Delta_b) = (f_1 B : \Delta_b B)/f_1 B$
- (ii) $\implies \mathfrak{a}B = (f_1 B : \Delta_b B)$
- Cramer's rule $\implies \Delta f_1, \Delta_b g_1 \in \mathfrak{b} \implies \Delta_b g_1 \equiv \Delta f_1 \pmod{\mathfrak{b}}$
 - $\implies (g_1 B : \Delta B) = (f_1 B : \Delta_b B) \implies (g_1 B : \Delta B) = \mathfrak{a}B$
 - $\implies \mathfrak{a}B/g_1 B = (g_1 B : \Delta B)/g_1 B = \text{Ann}_{B/g_1 B}((B/g_1 B)\Delta)$
- Canonical isomorphism $(f_1 B : \mathfrak{a}B)/f_1 B \cong (g_1 B : \mathfrak{a}B)/g_1 B$ sends image of Δ_b in $B/f_1 B$ to image of Δ in $B/g_1 B$
 - (i) $\implies (B/g_1 B)\Delta = (g_1 B : \mathfrak{a}B)/g_1 B = \text{Ann}_{B/g_1 B}(\mathfrak{a}B/g_1 B)$
- Conclude by noting $A' = B/g_1 B$, $\mathfrak{a}' = \mathfrak{a}B/g_1 B$, $\Delta' A' = \Delta(B/g_1 B)$

- Main Theorem $\implies \eta := \Theta^{-1}(1)$ generates $\text{Hom}_A(B, A)$
- $\Theta(\text{Tr}_{B/A}) = J = \mu(\Delta)$ and $\text{Tr}_{B/A} = J \cdot \eta$
- **Corollary:** Let A be reduced with $(x_1, \dots, x_n)^m \subseteq (t_1, \dots, t_n) \subseteq (x_1, \dots, x_n)$ for some m .
 - Write $t_j = \sum_{i=1}^n e_{ij}x_i$ and set $E \equiv \det(e_{ij}) \pmod{(t_1, \dots, t_n)}$. Then:
 - $J = (\dim_A B) \cdot E$
 - Get bilinear form $\beta_\eta : B \times B \rightarrow A$ defined by $\beta_\eta(x, y) = \eta(xy)$
 - Eisenbud-Levine: if $\phi : B \rightarrow A$ has $\phi(E) = 1$, then $\beta_\phi \cong \beta_\eta$
 - Kass-Wickelgren: β_η is the local \mathbb{A}^1 -degree

Questions?